

**Attack these protocols**

# Shamir-Rivest-Adleman three-pass

Goal: secrecy of  $M$ .

Assumed equality:  $\{\{x\}_k\}_{k'} \equiv \{\{x\}_{k'}\}_k$ .

1.  $A \rightarrow B: \{M\}_{K_a}$
2.  $B \rightarrow A: \{\{M\}_{K_a}\}_{K_b}$
3.  $A \rightarrow B: \{M\}_{K_b}$

(attacker finds  $M$ )

# Repeated authentication (several authors)

Goal:  $A$  and  $B$  make sure they are alive.

1.  $A \rightarrow B: N_a$
2.  $B \rightarrow A: N_b, \{N_a\}_{K_{ab}}$
3.  $A \rightarrow B: \{N_b\}_{K_{ab}}$

(attacker makes  $B$  believe  $A$  is alive without  $A$  doing anything))

# Woo-Lam (one-way authentication)

Goal:  $B$  makes sure that  $A$  is alive.

1.  $B \rightarrow A: N_b$
2.  $A \rightarrow B: \{N_b\}_{K_{as}}$
3.  $B \rightarrow S: B, \{A, \{N_b\}_{K_{as}}\}_{K_{bs}}$
4.  $S \rightarrow B: \{N_b\}_{K_{bs}}$

(attacker makes  $B$  believe  $A$  is alive without  $A$  doing anything)  
(how to repair?)

# Wide-mouthed frog

Goal: key exchange

1.  $A \rightarrow S: A, \{T_a, B, K_{ab}\}_{K_{as}}$
2.  $S \rightarrow B: \{T_s, A, K_{ab}\}_{K_{bs}}$

# Andrew Secure RPC

Goal: exchange of a new key  $K'_{ab}$

1.  $A \rightarrow B: A, \{N_a\}_{K_{ab}}$
2.  $B \rightarrow A: \{N_a + 1, N_b\}_{K_{ab}}$
3.  $A \rightarrow B: \{N_b + 1\}_{K_{ab}}$
4.  $B \rightarrow A: \{K'_{ab}, N'\}_{K_{ab}}$

# Yahalom

Goal: mutual authentication of  $A$  and  $B$ , key exchange

1.  $A \rightarrow B: A, N_a$
2.  $B \rightarrow S: B, N_b, \{A, N_a\}_{K_{bs}}$
3.  $S \rightarrow A: N_b, \{B, K_{ab}, N_a\}_{K_{as}}, \{A, K_{ab}, N_b\}_{K_{bs}}$
4.  $A \rightarrow B: \{A, K_{ab}, N_b\}_{K_{bs}}, \{N_b\}_{K_{ab}}$

(attacker impersonates Bob to Alice)

# Needham-Schroeder symmetric key

Goal: mutual authentication of  $A$  and  $B$ , key exchange

1.  $A \rightarrow S: A, B, N_a$
2.  $S \rightarrow A: \{N_a, B, K_{ab}, \{K_{ab}, A\}_{K_{bs}}\}_{K_{as}}$
3.  $A \rightarrow B: \{K_{ab}, A\}_{K_{bs}}$
4.  $B \rightarrow A: \{N_b\}_{K_{ab}}$
5.  $A \rightarrow B: \{N_b + 1\}_{K_{ab}}$

(attacker impersonates Alice to Bob)



# Kao-Chow Authentication

Goal: mutual authentication of  $A$  and  $B$ , key exchange

1.  $A \rightarrow S: A, B, N_a$
2.  $S \rightarrow B: \{A, B, N_a, K_{ab}\}_{K_{as}}, \{A, B, N_a, K_{ab}\}_{K_{bs}}$
3.  $B \rightarrow A: \{A, B, N_a, K_{ab}\}_{K_{as}}, \{N_a\}_{K_{ab}}, N_b$
4.  $A \rightarrow B: \{N_b\}_{K_{ab}}$

# Woo-Lam (key exchange)

Goal: mutual authentication of  $A$  and  $B$ , key exchange

1.  $A \rightarrow B: A, N_a$
2.  $B \rightarrow A: B, N_b$
3.  $A \rightarrow B: \{A, B, N_a, N_b\}_{K_{as}}$
4.  $B \rightarrow S: \{A, B, N_a, N_b\}_{K_{as}}, \{A, B, N_a, N_b\}_{K_{bs}}$
5.  $S \rightarrow B: \{B, N_a, N_b, K_{ab}\}_{K_{as}}, \{A, N_a, N_b, K_{ab}\}_{K_{bs}}$
6.  $B \rightarrow A: \{B, N_a, N_b, K_{ab}\}_{K_{as}}, \{N_a, N_b\}_{K_{ab}}$
7.  $A \rightarrow B: \{N_b\}_{K_{ab}}$

(attacker impersonates Alice to Bob, inserts own key)

# Needham-Schroeder public key

Goal: mutual authentication of  $A$  and  $B$

Assume public keys are bound to names using certificates or smth.

1.  $A \rightarrow B: \{[A, N_a]\}_{K_B}$
2.  $B \rightarrow A: \{[N_a, N_b]\}_{K_A}$
3.  $A \rightarrow B: \{[N_b]\}_{K_B}$